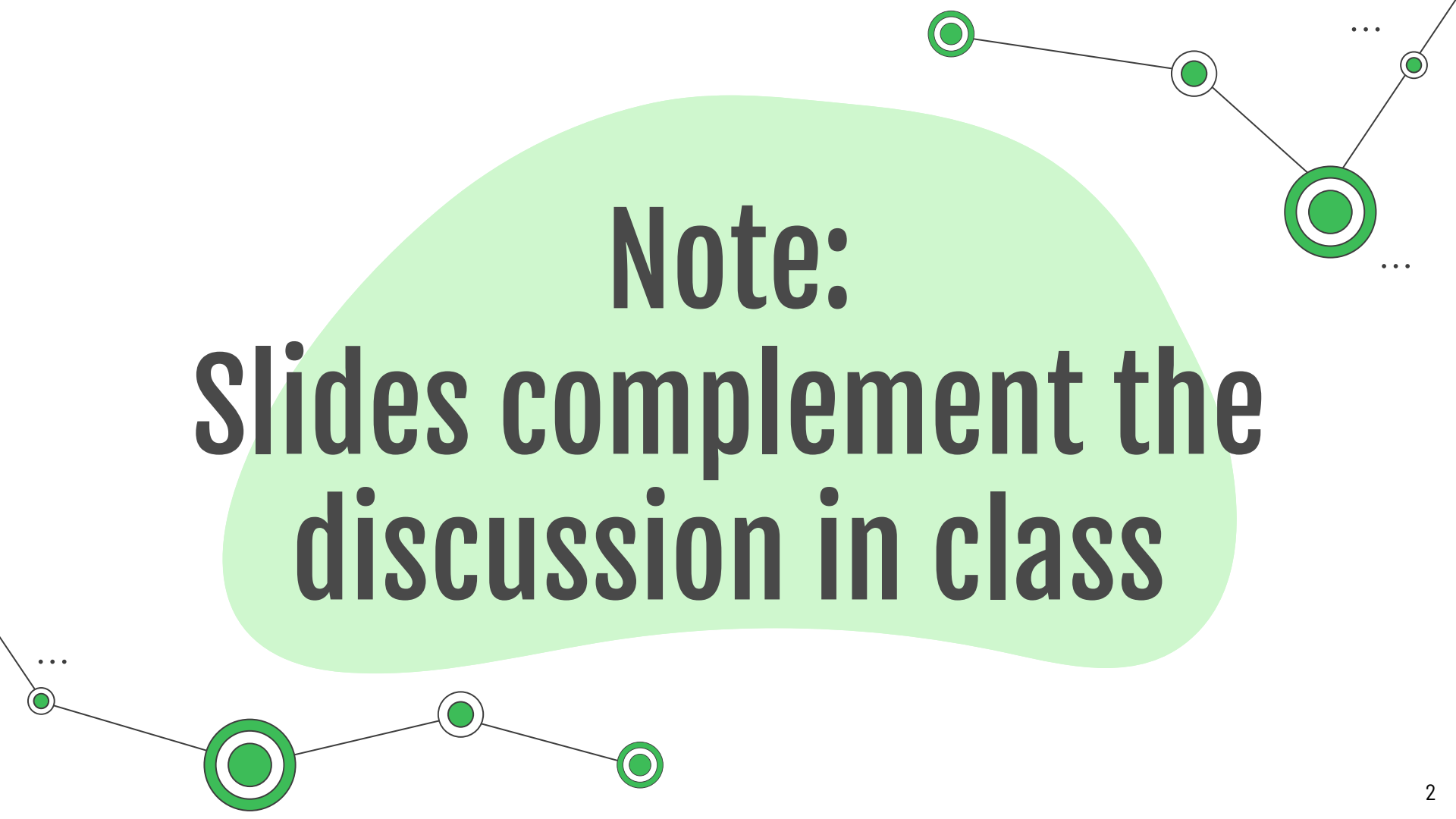


Array

CS 251 - Data Structures and
Algorithms

A decorative network diagram consisting of several green circular nodes connected by thin black lines. Some nodes are single green circles, while others are double green circles. The nodes are arranged in a non-linear fashion, with some at the top right and others at the bottom left. Ellipses (...) are placed near some nodes to indicate a larger, unseen network.

Note:
**Slides complement the
discussion in class**



Array

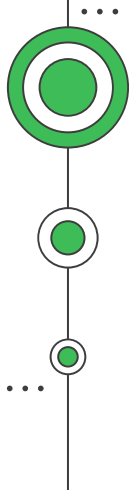
Static linear data structures

Table of Contents



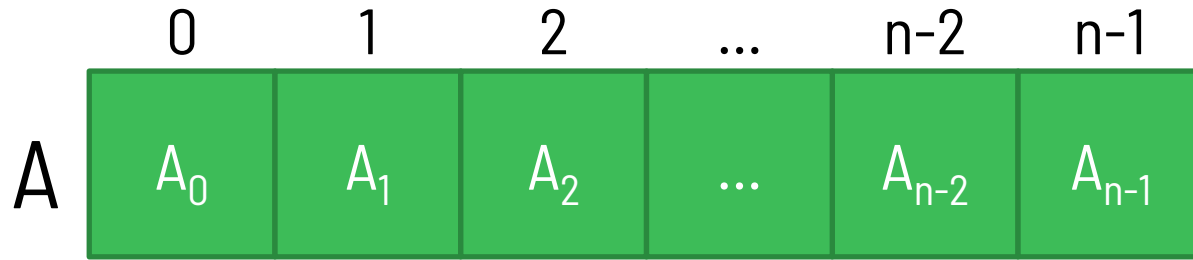
01 Array

Static linear data structure



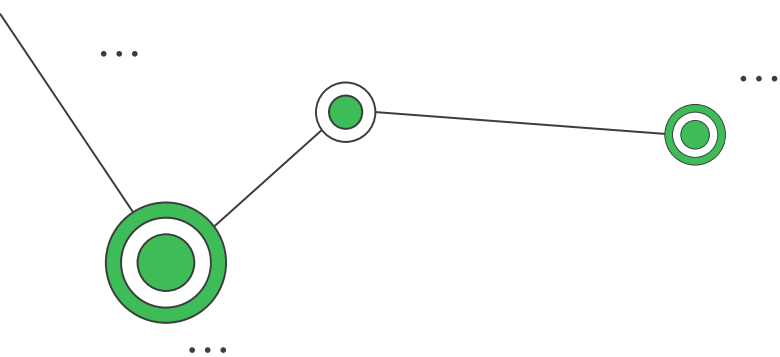


Array

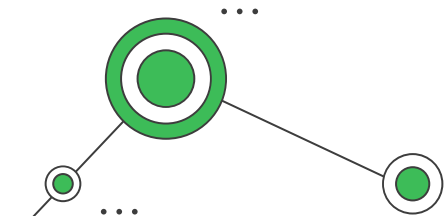


An array stores an item per index. An item could be as simple as a single value, or as complex as another data structure. Each item is accessed by its index in the array.

The size of an array is the number of items it stores. The capacity of an array is the amount of space reserved to store items. For example, an array could be of capacity 10 and size 3 (i.e., there are 3 items in it).

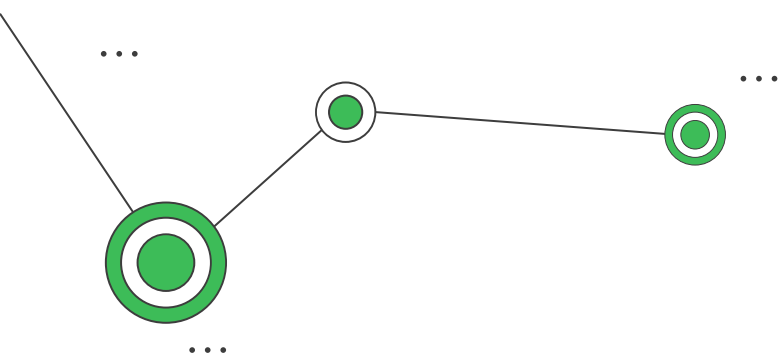


Insertion at the End

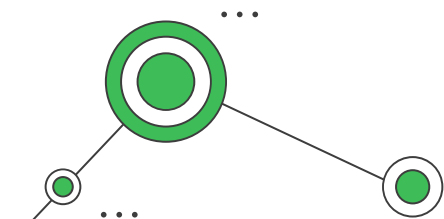


```
algorithm InsertEnd(A:array, x:item) → array
  let n be the size of A
  let m be the capacity of A
  if n = m then
    A ← Resize(A, 2 * m)
  end if
  A[n] ← x
  increase the size of A by 1
  return A
end algorithm
```

```
algorithm Resize(A:array, m: $\mathbb{Z}^+$ ) → array
  let B be an array of capacity m
  let n be the size of A
  for i from 0 to n-1 do
    B[i] ← A[i]
  end for
  set the size of B to n
  return B
end algorithm
```

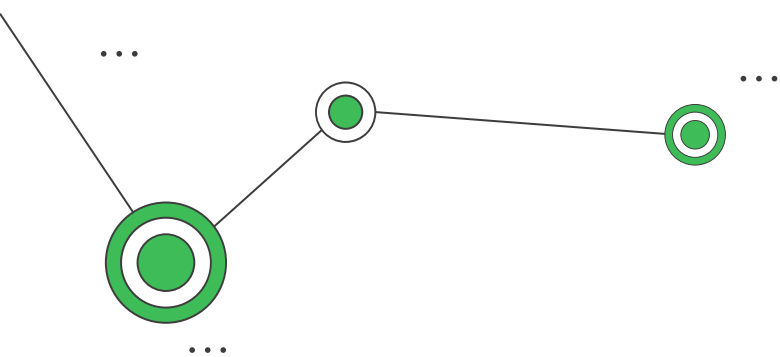


Insertion at Some Index



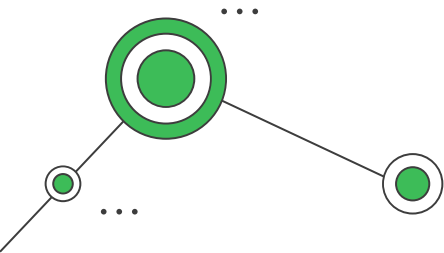
```
algorithm InsertAt(A:array, x:item, i: $\mathbb{Z}_{\geq 0}$ )  $\rightarrow$  array
  let n be the size of A
  let m be the capacity of A
  if n = m then
    A  $\leftarrow$  Resize(A, 2 * m)
  end if
  if i < n then
    A  $\leftarrow$  RightShift(A, i)
  end if
  A[i]  $\leftarrow$  x
  increase the size of A by 1
  return A
end algorithm
```

```
algorithm RightShift(A:array, i: $\mathbb{Z}_{\geq 0}$ )  $\rightarrow$  array
  let n be the size of A
  for j from n to i+1 by -1 do
    A[j]  $\leftarrow$  A[j-1]
  end for
  return A
end algorithm
```



Linear Search

```
algorithm LinearSearch(A:array, x:item) →  $\mathbb{Z}$ 
  let n be the size of A
  for i from 0 to n-1 do
    if A[i] = x then
      return i
    end if
  return -1
end algorithm
```





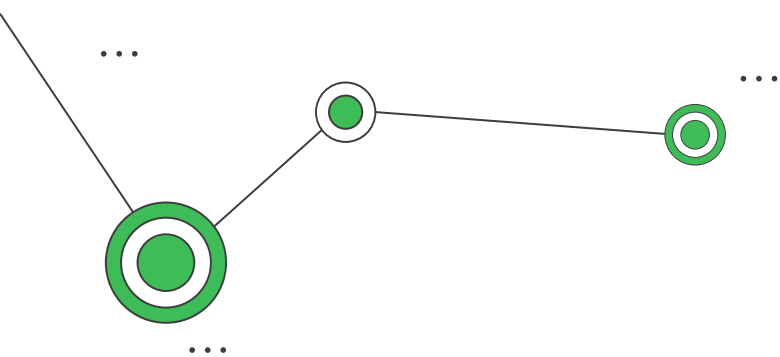
Recursive Linear Search



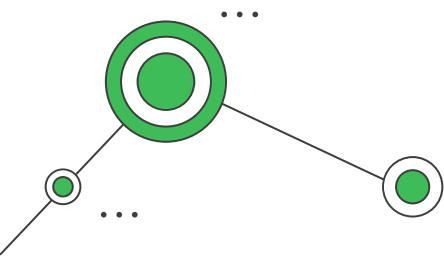
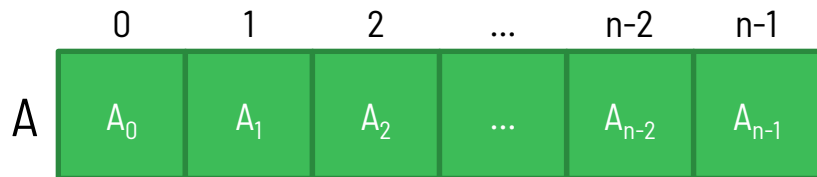
```
algorithm RLinearSearch(A:array, x:item, i: $\mathbb{Z}$ )  $\rightarrow \mathbb{Z}$ 
  if  $i < 0$  then
    return -1
  end if
  if  $A[i] = x$  then
    return i
  end if
  return RLinearSearch(A, x,  $i-1$ )
end algorithm
```

First call:

```
let n be the size of A
index  $\leftarrow$  RLinearSearch(A, x,  $n-1$ )
```



Array



Static Data Structure (i.e., fixed capacity)

Array access: $A[i] \in \Theta(1)$

Q: How many operations are required to **insert** an item into an unsorted array?

Keep track of next available cell?: $\Theta(1)$

Sorry, no tracking: $\Theta(n)$

Q: How much **space** is required to keep track of the next available cell?

One variable: $\Theta(1)$

Q: How many operations are required to **insert** an item into a sorted array?

Find location + Right shifting: $\Theta(n)$

Resize an Array



```
algorithm resize(A:array, m: $\mathbb{Z}^+$ )  $\rightarrow$  array  
  let n be the size of A  
  let B be an array of capacity m  
  copy/move the elements from A to B  
  return B or let A point to B?  
  perhaps delete A?  
end algorithm
```

Time? $T(n) \approx n \in \Theta(n)$

Space? $T(n) \approx n \in \Theta(n)$

Usual strategy:

Full? Increase its size (double the current size)

Half empty? reduce its size (half the current size)



2D Array (AKA Grid or Matrix)

A	0	1	2	...	m-2	m-1
0	$A_{0,0}$	$A_{0,1}$	$A_{0,2}$...	$A_{0,m-2}$	$A_{0,m-1}$
1	$A_{1,0}$	$A_{1,1}$	$A_{1,2}$...	$A_{1,m-2}$	$A_{1,m-1}$
2	$A_{2,0}$	$A_{2,1}$	$A_{2,2}$...	$A_{2,m-2}$	$A_{2,m-1}$
...
n-2	$A_{n-2,0}$	$A_{n-2,1}$	$A_{n-2,2}$...	$A_{n-2,m-2}$	$A_{n-2,m-1}$
n-1	$A_{n-1,0}$	$A_{n-1,1}$	$A_{n-1,2}$...	$A_{n-1,m-2}$	$A_{n-1,m-1}$

Static Data Structure (i.e., fixed size)

Array access: $A[i][j] \in \Theta(1)$

Q: How many operations are required to traverse a grid?

General case: $\Theta(nm)$

Squared matrix: $\Theta(n^2)$

Q: How much space is required to store a grid?

General case: $\Theta(nm)$

Squared matrix: $\Theta(n^2)$



EOF

Do you have any questions?

CREDITS: This presentation template was created by [Slidesgo](#), including icons by [Flaticon](#), infographics & images by [Freepik](#) and illustrations by [Stories](#)